

# Efficiently Computing LS

Monday, April 10, 2017 9:03 AM

## Summary of Least squares problems

pseudo-inverse when  $A$  is full column rank

$$\hat{x} = (A^T A)^{-1} A^T y$$

pseudo-inverse when  $A$  is full row rank

$$\hat{x} = A^T (A A^T)^{-1} y$$

Tikhonov regularization

$$\hat{x} = (A^T A + \delta I)^{-1} A^T y$$

generalized Tikhonov reg.

$$\hat{x} = (A^T A + \delta D^T D)^{-1} A^T y$$

weighted least squares

$$\hat{x} = (A^T W^T W A)^{-1} A^T W^T W y$$

BLUE

$$\hat{x} = (A^T R^{-1} A)^{-1} A^T R^{-1} y$$

each solves a symmetric, positive definite set of equations

There are better ways to solve these than with direct computation

if we have a triangular system

$$\begin{bmatrix} a_{11} & 0 & & \\ a_{21} & a_{22} & & \\ \vdots & & \ddots & \\ a_{n1} & & & a_{nn} \end{bmatrix} x = y$$

(solve for  $x$ )

$$x_1 = y_1 / a_{11}$$

$$x_2 = \frac{1}{a_{22}} (y_2 - a_{21} \cdot x_1)$$

$\vdots$

we can solve this with  $O(N^2)$  operations instead of

$$x = A^{-1} y \leftarrow O(N^3)$$

we know each of these terms

# Exploiting Structure

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if  $A\hat{x} = y$   
└ solving for  $x$

if  $A$  is diagonal  $O(N)$

if  $A$  is triangular  $O(N^2)$ \*

if  $A$  is orthogonal  $O(N^2)$ \*

$$\hat{x} = A^T y$$

some clever algorithms

Levinson-Durbin algorithm

if  $A$  is Toeplitz  $O(N^2)$  correlation matrices

if  $A$  is Circulant  $O(N \log N)$

if  $A$  is general  $O(N^2)$

$$\begin{pmatrix} 1 & 2 & 3 & 4 \\ 10 & 1 & 2 & 3 \\ 7 & 10 & 1 & 2 \\ 3 & 7 & 10 & 1 \end{pmatrix}$$

Performing a Cholesky decomposition on sym+def  $A$

$A = LL^T$  ← we can use the triangular method twice to find a solution to  $Ax = y$   $O(N^2)$

same for QR decomposition  
(can be found using Gram-Schmidt)

$A = QR$   $O(N^2)$   
└ upper triangular  
└ orthonormal

used in communications

often used if the  $Ax = y$  equation needs to be solved for multiple  $y$ 's. But, they are often used even for a single  $y$  because they are more numerically stable than solutions involving  $A^{-1}$ .

# Iterative methods for Least Squares

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$$\hat{x} = \underbrace{(A^T A)^{-1}}_{\text{costs } O(MN^2)} A^T y \quad A \in \mathbb{R}^{m \times n}$$

$$A^T A \hat{x} = A^T y$$

solving this system directly is  $O(N^3)$

example: MRI  $M = 5 \cdot 10^6$  by  $N = 2 \cdot 10^6$   
 $MN^2 \sim 10^{19}$

for such a system, we need another solution

Recasting as an optimization

$$\underbrace{A^T A} x = \underbrace{A^T y}$$
$$H x = b$$

(sym+def)

The following is a quadratic program and is convex

$$\min_{x \in \mathbb{R}^n} \frac{1}{2} x^T H x - x^T b$$

proceed by setting the derivative to zero

$$\nabla_x \left( \frac{1}{2} x^T H x - x^T b \right) \Big|_{x=\hat{x}} = 0$$

$\Downarrow$

$$H x - b = 0$$

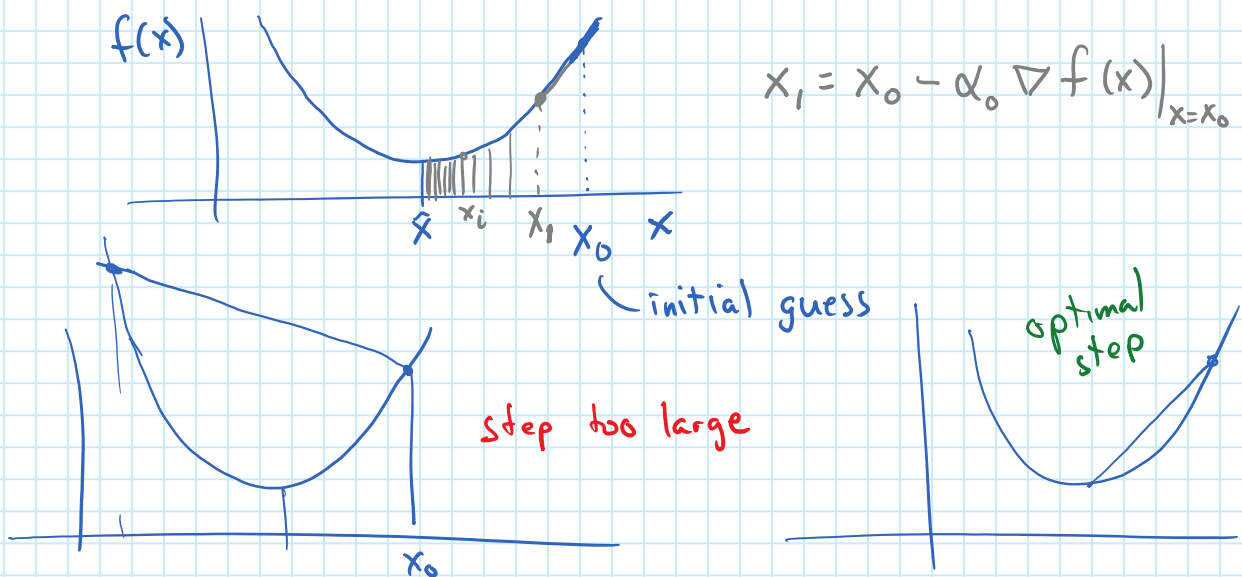
$\Downarrow$

$$H x = b$$

# Gradient Descent

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$$f(x) = x^T H x - x b \quad \leftarrow \text{one dimensional example}$$



we can define  $\frac{b - H x_k}{-\nabla f(x) \Big|_{x=x_k}} = r_k$  residual

$$x_{k+1} = x_k + \alpha_k r_k \quad \leftarrow \text{update equation}$$

to make  $f(x_{k+1})$  as small as possible

$$f(x_k + \alpha_k r_k) \quad \leftarrow \text{we can solve for } \alpha \text{ that minimizes this}$$

$$\frac{d}{d\alpha} f(x_k + \alpha r_k) = 0$$

$$= \nabla f(x_{k+1})^T \frac{d}{d\alpha} x_{k+1} \quad \leftarrow \text{using chain rule}$$

$$= \nabla f(x_{k+1})^T r_k$$

we need  $r_k \perp \nabla f(x_{k+1})$

$$\text{Thus } r_{k+1} \perp r_k$$

# Gradient Continued

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$$\begin{aligned}
 r_{k+1}^T r_k &= 0 \\
 (b - Hx_{k+1})^T r_k &= 0 \\
 (b - H(x_k + \alpha_k r_k))^T r_k &= 0 \\
 \vdots \\
 \alpha_k &= \frac{r_k^T r_k}{r_k^T H r_k}
 \end{aligned}$$

ensures rapid convergence.  
 ↓  
 for a large problem  
 20-50 steps.

In general, moving in the proper direction with a non-zero, but sufficiently small step size will yield convergence.

## LMS algorithm

we are trying to minimize  $E\{e^2[n]\}$  where

$$e[n] = d[n] - \underline{w}^T \underline{x}[n]$$

$$\underline{x}[n] = \begin{pmatrix} x[n] \\ x[n-1] \\ \vdots \\ x[n-p] \end{pmatrix}$$



$$\nabla_w E\{e^2[n]\} = E\{\nabla_w e^2[n]\} = E\{2e[n] \nabla_w e[n]\}$$

$$= -E\{2e[n] \underline{x}[n]\} = 0$$

gives the solution exactly

# LMS Algorithm

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$$E\{e[n] \underline{x}[n]\} = E\{(d[n] - \underbrace{\underline{w}^T \underline{x}[n]}_{\text{scalar}}) \underline{x}[n]\}$$

$$= E\{d[n] \underline{x}[n]\} - E\{\underline{x}[n] \underline{x}^T[n]\} \underline{w}$$

cross correlation of  $\underline{x}[n]$  with  $d[n]$

The normal equation

$$= \underline{r}_x - R_x \underline{w} = 0 \Rightarrow R_x \underline{w} = \underline{r}_x$$

we have seen problems like this before!

instead of solving the normal equation directly, we will use gradient descent

$$\underline{w}_{n+1} = \underline{w}_n - \frac{1}{2} \mu \nabla_{\underline{w}} E\{e^2[n]\}$$

$$\underline{w}_{n+1} = \underline{w}_n + \mu E\{e[n] \underline{x}[n]\} \quad \text{from ①}$$

how to find?

$$\hat{E}\{e[n] \underline{x}[n]\} = \frac{1}{N} \sum_{k=n-N+1}^n e(k) \underline{x}(k)$$

(sample mean)

The LMS algorithm uses

$$E\{e[n] \underline{x}[n]\} \approx e[n] \underline{x}[n]$$

$$\underline{w}_{n+1} = \underline{w}_n + \mu e[n] \underline{x}[n]$$

# LMS Algorithm Practical Issues

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$$y[n] = \underline{w}^T[n] \underline{x}(n) \quad \underline{x}(n) = \begin{pmatrix} x(n) \\ x(n-1) \\ \vdots \\ x(n-p) \end{pmatrix} \quad \underline{w}[n] = \begin{pmatrix} w_0(n) \\ w_1(n) \\ \vdots \\ w_p(n) \end{pmatrix}$$

$$e(n) = d(n) - y(n)$$

↑ "d" for desired or target signal

$$\underline{w}[n+1] = \underline{w}[n] + \mu e(n) \underline{x}[n] \quad \text{minimizes the expected error, } E\{e^2(n)\}$$

$$R_x = E\{\underline{x}[n] \underline{x}^T[n]\} \quad \leftarrow \text{correlation matrix}$$

if  $\{\lambda_i\}$  are the eigenvalues of  $R_x$  then

$$0 < \mu < \frac{2}{\lambda_1} \quad \leftarrow \lambda_{\max}$$

$$\lambda_{\max} \leq \sum_{k=0}^p \lambda_k = \text{Tr}(R_x) = \sum_{k=0}^p E\{x^2(n-k)\} = \dots$$

if  $x(n)$  is stationary (or WSS)

$$= (p+1) \cdot (E\{x^2(n)\}) \approx \sum_{k=0}^p x^2(n-k) = \underline{x}^T(n) \underline{x}(n)$$

approx result from assumption in convergence proof.

$$0 < \mu < \frac{2}{\lambda_{\max}} \Rightarrow 0 < \mu < \frac{2}{\underline{x}^T(n) \underline{x}(n)}$$

we define  $\mu(n) = \frac{\beta}{\underline{x}^T(n) \underline{x}(n)}$  where  $0 < \beta < 2$

in practice, we choose  $\beta = 0.5$   
(or in the range 0.1 to 1)

# nLMS Adaptive Filters

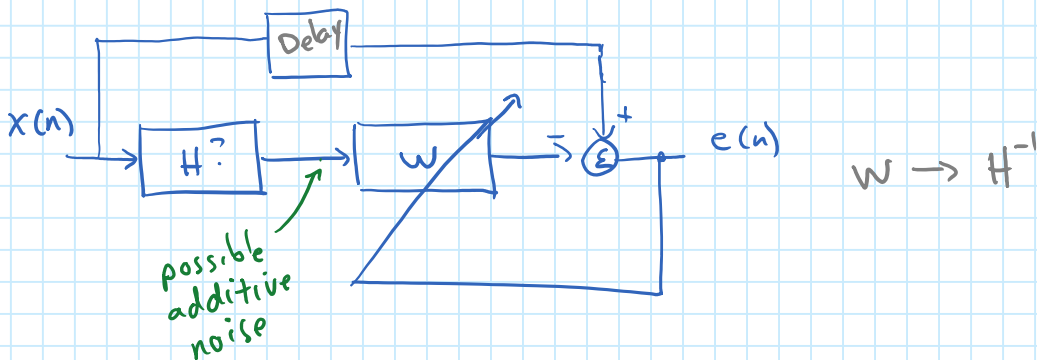
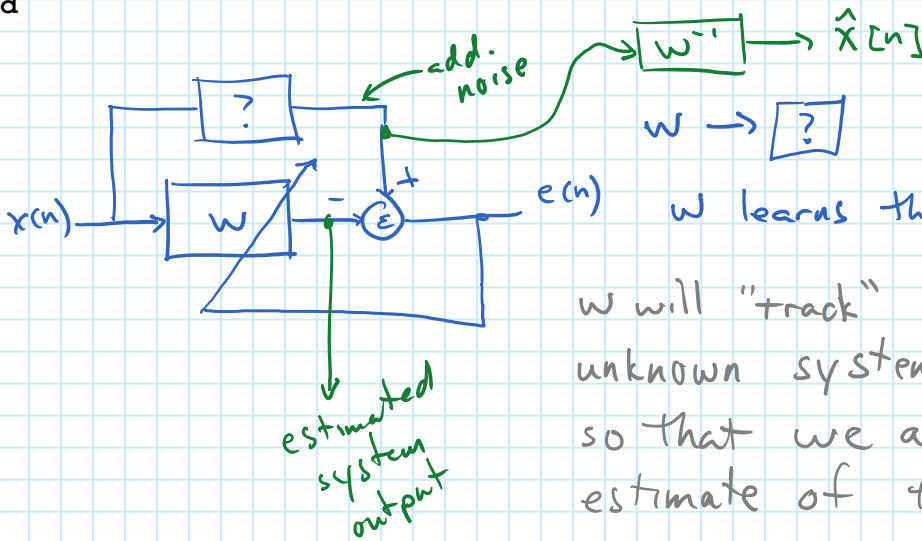
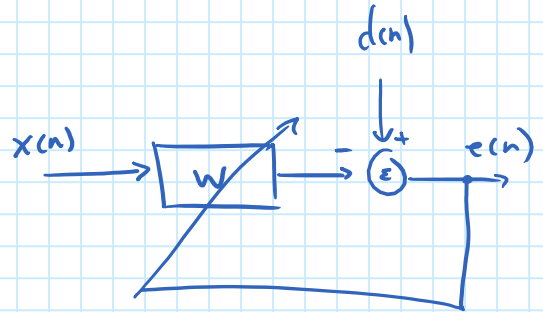
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$$\underline{w}[n+1] = \underline{w}[n] + \frac{\beta}{\|\underline{x}(n)\|^2} \underline{x}(n) e(n)$$

```
function [y,e,w]=mylms(x,d,w,beta)
% length x = length d

x=x(:); d=d(:); w=w(:);
y = zeros(size(x));
e = zeros(size(x));

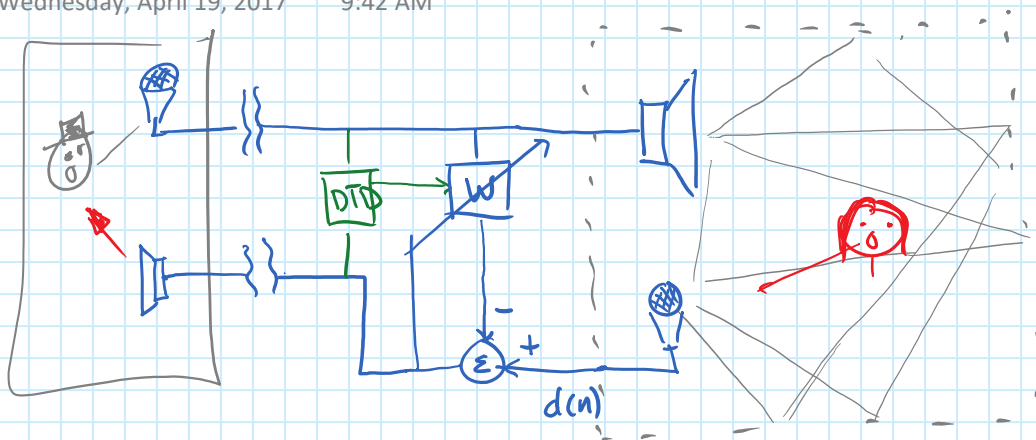
for i=length(w):length(x),
    xx = x(i:-1:i-length(w)+1);
    y(i) = w' * xx;
    e(i) = d(i) - y(i);
    w = w + beta/(xx'*xx+1e-8) * e(i) * xx;
end
```





# Echo Cancellation

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Adaptive LMS filters adapt most quickly when  $x[n]$  is white noise.

↳ adaptation rate is a function of  $\frac{\lambda_{\max}}{\lambda_{\min}}$  → larger spread = slower adaptation